

# CTEC323 Lecture 8

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# What is Normalization?

- ▶ Remember all of the data redundancy problems with using files from Chapter 1 (update, insert, and delete anomalies)?
- ▶ Without good table structure, databases can have these problems, too!
- ▶ Normalization: process for evaluating and correcting table structures to minimize data redundancies, thereby reducing the likelihood of data anomalies
- ▶ The higher the normal form, the more join operations required to produce a specified output
- ▶ Sometimes, denormalizing a database is required to meet performance requirements (but, of course, this increases data redundancy)

# Running Example

We begin with a table with the following attributes.

proj_num	proj_name	emp_num	emp_name
job_class	chg_hour	hours	

- ▶ Assume each project is completed by many employees, and each employee can work on multiple projects
- ▶ Each employee has exactly one job\_class and each job\_class has one chg\_hour
- ▶ hours represents the hours worked on a particular project by a particular employee

# Problems with this Table

- ▶ This structure invites data inconsistencies
  - ▶ What if we enter Elect. Engineer for the job\_class of one employee and Elect. Eng for others?
- ▶ It also promotes anomalies, such as
  - ▶ Updating an employee's job\_class may require changes in many rows
  - ▶ Updating a job's chg\_hour may require changes in many rows
  - ▶ If we insert a new employee who has no project yet, we have to fill in the project information with phantom data
  - ▶ If we delete the sole employee of a project, we lose the project information too

# Objectives of Normalization

- ▶ Produce tables that represent single subjects (courses, students, etc.),
- ▶ Avoid storing data unnecessarily in more than one table
- ▶ Ensure that attributes in a table are dependent on
  - ▶ the primary key (1NF)
  - ▶ the entire primary key (2NF)
  - ▶ nothing but the primary key (3NF)

# 1NF

- ▶ All relational tables are already in 1NF (first normal form)
- ▶ All we have to do is identify a PK
- ▶ The combination of proj\_num and emp\_num gives us a PK for our running example

# Identifying Dependencies

1.  $\text{proj\_num, emp\_num} \rightarrow \text{proj\_name, emp\_name, job\_class, chg\_hour, hours}$
  2.  $\text{proj\_num} \rightarrow \text{proj\_name}$
  3.  $\text{emp\_num} \rightarrow \text{emp\_name, job\_class, chg\_hour}$
  4.  $\text{job\_class} \rightarrow \text{chg\_hour}$
- ▶ Partial dependency: dependency based on only a part of a PK (dependencies 2 and 3 above)
  - ▶ Transitive dependency: dependency of one non-PK attribute on another non-PK attribute (dependency 4 above)
  - ▶ If we start with a 1NF table and remove the partial dependencies, we have a 2NF table

## 2NF

- ▶ To convert from 1NF to 2NF, we take each component of the PK that functionally determines other attributes, copy that part of the PK to a new table, and move the attributes it determines from the original table to that new table
- ▶ Each component of the original key becomes the PK of its new table
- ▶ For our example, this results in three tables (PK attributes in bold)
- ▶ project (**proj\_num**, proj\_name)
- ▶ employee (**emp\_num**, emp\_name, job\_class, chg\_hour)
- ▶ assignment (**proj\_num**, **emp\_num**, hours)
- ▶ Note: this format of writing tables with PK's in bold is called a relational schema

# 3NF

- ▶ A table in 2NF may still have transitive dependencies; we eliminate these when we convert a 2NF table into 3NF
- ▶ To convert a table from 2NF to 3NF, begin by identifying each determinant of a transitive dependency and the attributes it determines
- ▶ Then, we take each determinant, copy it to a new table, and move the attributes it determines from the original table to that new table
- ▶ Each determinant becomes the PK of its new table

## 3NF...

- ▶ The project and assignment tables have no determinants, so they are already in 3NF
- ▶ Employee has a determinant:  $\text{job\_class} \rightarrow \text{chg\_hour}$
- ▶ Converting to 3NF, we split employee into two tables
- ▶ employee (**emp\_num**, emp\_name, job\_class)
- ▶ job (**job\_class**, chg\_hour)
- ▶ Our final database with all tables in 3NF thus contains four tables

## 3NF...

- ▶ assignment (**proj\_num**, **emp\_num**, hours)
- ▶ project (**proj\_num**, proj\_name)
- ▶ employee (**emp\_num**, emp\_name, job\_class)
- ▶ job (**job\_class**, chg\_hour)

## Numeric Keys

- ▶ Using job\_class as the PK in the employee table forces us to enter the text for the job class (such as Database Designer) for each new employee
- ▶ If we make a spelling mistake on the job\_class of a new employee, it will not match a job\_class in job (referential integrity problem)
- ▶ An alternative is to add a job\_code attribute which becomes the new PK of the job table (and a new FK in the employee table)
- ▶ employee (**emp\_num**, emp\_name, job\_code)
- ▶ job (**job\_code**, job\_class, chg\_hour)
- ▶ This introduces a transitive dependency (which we could eliminate)
- ▶ job\_code is called a surrogate key because it is an artificial key introduced to simplify the assignment of PK's to tables

# Historical Accuracy

- ▶ We can add the job charge per hour to the assignment table to capture the amount charged per hour at the time the project was completed
- ▶ If we don't store this value and the assignment table and the job charge later changes, we will incorrectly calculate the total charge of a project (we will not use the historically accurate value for the hourly rates)
- ▶ assignment (**proj\_num**, **emp\_num**, hours, assign\_chg\_hour)

# Denormalization

- ▶ As tables are decomposed in the normalization process, the number of tables expands
- ▶ Retrieving meaningful information often requires multiple joins and therefore increased system resources
- ▶ We occasionally denormalize to increase processing speed
- ▶ Unnormalized tables have the following drawbacks
  - ▶ Greater possibility for data anomalies
  - ▶ Updates are less efficient because larger tables have to be processed